## Graph Theory Homework 3

## Joshua Ruiter

## January 27, 2018

**Exercise 1.** Let G = (V, E) be a bipartite graph,  $V = A \sqcup B$ . Let  $\widetilde{G}$  be the network formed from G by vertices s, t, and directed edges  $\{(s, a) : a \in A\} \cup \{(b, t) : b \in B\}$ , and changing all edges in in G to directed edges in  $\widetilde{G}$  pointing from A to B. We give each edge in  $\widetilde{G}$  capacity 1. Finding a maximal flow in  $\widetilde{G}$  is equivalent to finding a maximal matching in G.

Augmenting Path Algorithm	Maximal Matching Algorithm
1. Initialize a flow $f_0$ with value zero	1. Define $M_0 = \emptyset$ to be our initial collection
everywhere.	of edges in a matching.
2. Choose a below-capacity path from	2. Choose an edge $e \in E$ that doesn't use any
s to t and set $f_{k+1} = f_k + 1$ on each	vertex used by an edge in $M_k$ , and set
edge of that path.	$M_{k+1} = M_k \cup \{e\}.$
3. Augment $f_k$ using a backward path	3. If there is an edge $(a_1, b_1) \in M_k$ and edges
if possible.	$(a_1, b_2), (a_2, b_1) \in E$ , with $a_2, b_2$ not already used
	by edges in $M_k$ , then set
	$M_{k+1} = (M_k \setminus \{(a_1, b_1)\}) \cup \{(a_1, b_2), (a_2, b_1)\}$
4. Repeat steps 2-3 until no more	4. Repeat steps 2-3 until no more
augmenting is possible.	augmenting is possible.

The picture for step 3 of the Maximal Matching Algorithm looks like this. The diagonal edge is already in  $M_k$ , and the horizontal edges are the ones we replace it with.



**Proposition 0.1** (Exercise 2ab). Let G be a finite group and H a subgroup. Then there is a set  $\{g_1, \ldots, g_k\}$  of representatives for the left cosets of H that is simultaneously a set of representatives for the right cosets of H if and only if any collection of n distinct left cosets intersects at least n distinct right cosets.

*Proof.* We form a bipartite graph  $\widetilde{G}$  that contains information about cosets of G and apply the Hall Matching Theorem. The vertices of  $\widetilde{G}$  are the collection of all left and right cosets

of H in G. There is an edge between two cosets  $g_i H$  and  $Hx_j$  if  $g_i H \cap Hx_j \neq \emptyset$ .

$$V = L \sqcup R = \{g_1H, \dots, g_mH\} \sqcup \{Hx_1, \dots, Hx_m\}$$
$$E = \{(g_iH, Hx_j) : g_iH \cap Hx_j \neq \emptyset\}$$

Given a complete matching in  $\tilde{G}$  from L to R, we get simultaneous coset representatives, since each left coset has exactly one edge to a right coset, and this edge represents a simultaneous representative for those two cosets.

By Hall's Matching Theorem, a complete matching on this graph exists if and only if any set of n vertices in L has at least n neighbors in R. A collection of left cosets (vertices in L) has a neighbor for each right coset that intersects the union of that collection. That is, there is a complete matching if and only if for a collection of n distinct left cosets, their union intersects at least n distinct right cosets.

**Proposition 0.2** (Exercise 2c). Let G be a group, and  $H \subset G$  a finite subgroup. If  $\{g_1H, \ldots, g_nH\}$  is a collection of n distinct left cosets of H in G, then  $\bigcup_i g_iH$  intersects at least n distinct right cosets of H in G.

*Proof.* Let  $\{Hx_1, \ldots, Hx_k\}$  be the collection of (distinct) right cosets that intersect  $\bigcup_i g_i H$ . Since the right cosets cover all of G, this implies

$$\bigcup_{i=1}^n g_i H \subset \bigcup_{i=1}^k H x_i$$

Since left cosets are pairwise disjoint,  $\bigcup_i g_i H$  has n|H| elements. Since right cosets are pairwise disjoint,  $\bigcup_i Hx_i$  has k|H| elements. Thus

$$\bigcup_{i=1}^{j} g_i H \subset \bigcup_{i=1}^{k} H x_i \implies n|H| \le k|H| \implies n \le k$$

Thus  $\bigcup_i g_i H$  intersects at least n disjinct right cosets.

Interestingly, the previous proposition doesn't require G to be finite. However, the equivalence in Proposition 0.1 does need G to be finite, since Hall's Matching Theorem only includes finite graphs (as far as I know).

**Corollary 0.3** (Exercise 2c). If G is a finite group and H is a subgroup, there exists a set  $\{g_1, \ldots, g_k\}$  of representatives for the left cosets of H that is simultaneously a set of representatives for the right cosets of H.

*Proof.* By Proposition 0.2, the second half of the equivalence in Proposition 0.1 always holds.  $\Box$